NIST 14

brief comments on some of the 14 Round 2 NIST SHA-3 hash function candidates

Stanisław Radziszowski Rochester Institute of Technology spr@cs.rit.edu

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Block cipher based designs

AES based

- ECHO, Gilbert+ FR, wide-pipe
- Fugue, Jutla+ IBM, sponge-like
- Grøstl, Knudsen+ DK, wide-pipe
- SHAvite-3, Dunkelman+ IS, narrow-pipe

Other block cipher based

- BLAKE, Aumasson+ CH, narrow-pipe, ARX
- Hamsi (+-), Küçük BE/TR, narrow-pipe, bitsliced
- Skein, Schneier+ US, narrow-pipe, ARX, Threefish



Keccak The Keccak Team, BE

Team (from STMicroelectronics and NXP Semiconductors): Guido Bertoni, Joan Daemen (of AES fame), Michaël Peeters, Gilles Van Assche

- Elegant, convincing design, ideas from Grindahl
- Runs on a $5 \times 5 \times 2^{l}$ cube of bits, recommended 1600-bit state (l = 6)
- Nontrivial padding, message schedule
- Same rounds all over, except constants RC
- Sponge construction

This is my favorite hash!



Keccak single round, only one nonlinear step χ

$$\begin{split} \mathbf{R} &= \iota \circ \chi \circ \pi \circ \rho \circ \theta, \text{ with} \\ \theta: \ a[x][y][z] \ \leftarrow a[x][y][z] + \sum_{y'=0}^{4} a[x-1][y'][z] + \sum_{y'=0}^{4} a[x+1][y'][z-1], \\ \rho: \ a[x][y][z] \ \leftarrow a[x][y][z-(t+1)(t+2)/2], \\ & \text{ with } t \text{ satisfying } 0 \leq t < 24 \text{ and } \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix}^{t} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} x \\ y \end{pmatrix} \text{ in } \mathrm{GF}(5)^{2 \times 2} \\ & \text{ or } t = -1 \text{ if } x = y = 0, \\ \pi: \ a[x][y] \ \leftarrow a[x'][y'], \text{ with } \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 2 & 3 \end{pmatrix} \begin{pmatrix} x' \\ y' \end{pmatrix}, \\ \chi: \ a[x] \ \leftarrow a[x] + (a[x+1]+1)a[x+2], \\ \iota: \ a \ \leftarrow a + \mathrm{RC}[i_{\mathrm{r}}]. \end{split}$$



Skein

Schneier-Ferguson + team, US-UK

The main designers are well known authors of cryptography and security books

- Based on block cipher Threefish (Twofish competed in 2000 to become AES)
- Separate slides by John Hicks with a look inside
- Tree hashing strongly stressed
- Easy parallelizability

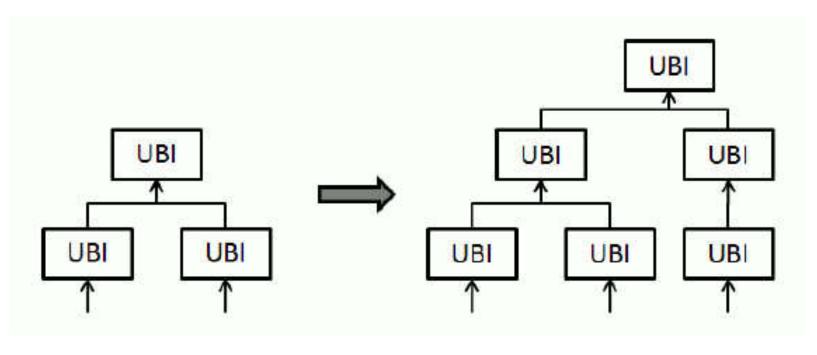
The only candidate among remaining 14 with easy and natural parallelizability!





tree hashing mode, tree growth

Unique Block Iteration - UBI

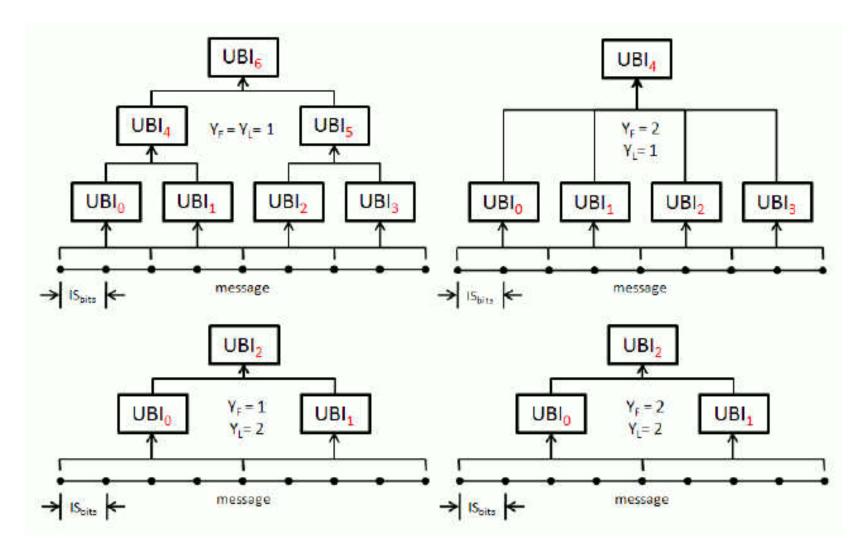


[Schorr 2010]



Skein

tree hashing mode, various parameters

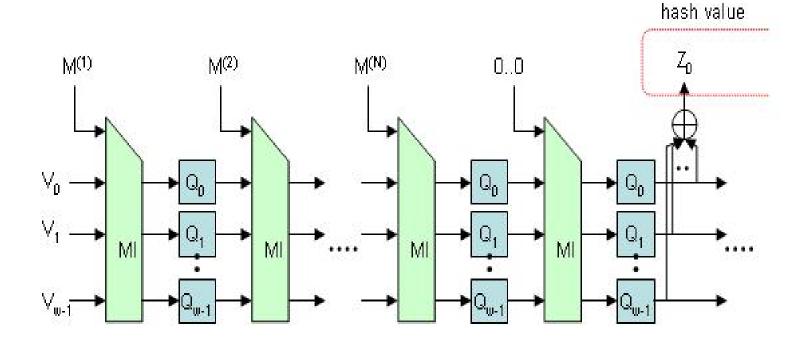


[Schorr 2010]



Luffa

Watanabe-Sato-DeCarniere, JP/BE



- MI linear maps, Q perms, sponge-like
- Small s-boxes motivated by Serpent (almost AES to be 10 years ago)
- Black horse?



Strong among AES based candidates

- Detailed slides by Joel Lathrop
- Wide-pipe Merkle-Damgård design
- Wide-trail design (not that clear how different from the latter)
- (512 or 1024)-bit blocks and chaining values
- (10 or 14) AES-type rounds on 8 by (8 or 16) array of bytes
- AES S-box



Hamsi Özgül Küçük, KU Leuven, BE

hamsi = sardela = anchovy

- Small message blocks, 32 and 64 bits
- Internal state size 512 and 1024 bits
- Message expansion via linear codes over GF(4) 32 to 256 bits via [128,16,70] code 64 to 512 bits via [256,32,131] code provides strong diffusion
- Sponge-like construction
- Small S-box from Serpent
- XOR of constants and a counter



CubeHash *r/b – h*

Daniel Bernstein, Chicago

• *r* rounds per each *b*-byte block, *h*-bit hash recommended r = 16, b = 32, sponge construction

new tweak on initialization/finalization with less rounds

• Nicely coded with 5-dimensional (cube) arrays

x[2][2][2][2][2]=y[32]

a round has ten simple steps manipulating

 $x_{ijklm} = x[i][j][k][l][m]$

with addition modulo 2³², word rotation, swap and XOR.

 Subject of vigorous discussions on the NIST hash-forum, perhaps the most studied among all candidates



BMW - Blue Midnight Wish

team from NTNU Trondheim, NO

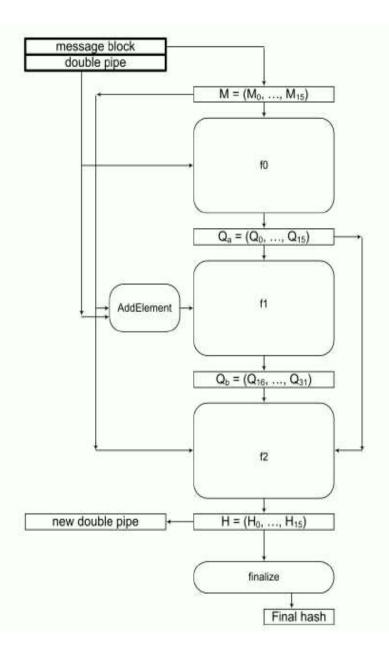
Danilo! Gligoroski, Vlastimil Klima, Svein Johan Knapskog and the team from the Norwegian University of Science and Technology

- Double-pipe expansion
- Very complicated details of specification
- Compression: non-linear expansion followed by folding



BMW - Blue Midnight Wish

structure





References

containing pointers to many other references

- NIST report NISTIR 7620, Status Report on the First Round of the SHA-3 Cryptographic Hash Algorithm Competition, September 2009 http://csrc.nist.gov/groups/ST/hash
- The Second SHA-3 Candidate Conference, Santa Barbara, CA, Aug. 23-24, 2010
 http://csrc.nist.gov/groups/ST/hash/ sha-3/Round2/Aug2010
- ECRYPT SHA-3 Zoo

http://ehash.iaik.tugraz.at/wiki/The_SHA-3_Zoo

• eBASH: ECRYPT Benchmarking of All Submitted Hashes http://bench.cr.yp.to/ebash.html

