Rabin Karp

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Overview Of Presentation

- What it is?
- How it works?
- Reduced Calculation
- Algorithm and Example
- Complexity
- Applications

What Rabin Karp Algorithm is?

- It is String matching algorithm.
- It is another application of Hashing.
- It is widely used for multiple pattern search.

Concept of Rabin Karp Algorithm

- The **Rabin-Karp** string searching algorithm calculates a **hash value** for the pattern, and for each M-character subsequence of text to be compared.
- If the hash values are unequal, the algorithm will calculate the hash value for next M-character sequence.
- If the hash values are equal, the algorithm will compare the pattern and the M-character sequence.
- In this way, there is only one comparison per text subsequence, and character matching is only needed when hash values match.

Some Questions for R.K.

- What is the hash function used to calculate values for character sequences?
- Isn't it time consuming to hash every one of the M-character sequences in the text body?
- To answer these question we refer to some mathematics.

Some Mathematics for R.K.

- Consider an M-character sequence as an M-digit number in base ***b***, where ***b*** is the number of letters in the alphabet. The text subsequence t[i .. i+M-1] is mapped to the number

  $x(i) = t[i]*b^{M-1} + t[i+1]*b^{M-2} + \ldots + t[i+M-1]$
- Furthermore, given x(i) we can compute x(i+1) for the next subsequence t[i+1 .. i+M] in constant time,as follows:

  $x(i+1) = t[i+1]*b^{M-1} + t[i+2]*b^{M-2} + \ldots + t[i+M]$

Mathematics Continue

$x(i+1) = x(i)*b$ (Shift left one digit)

$- t[i]*b\text{^}M$ (Subtract leftmost digit)

$+ t[i+M]$ Add new rightmost digit

- In this way, we never explicitly compute a new value. We simply adjust the existing value as we move over one character.
- If M is large, then the resulting value (b^M) will be enormous. For this reason, we hash the value by taking it **mod** a prime number q

Some more Mathematics…

- The **mod** function is particularly useful in this case due to several of its inherent properties:-

$[(x \bmod q) + (y \bmod q)] \bmod q = (x+y) \bmod q$

$(x \bmod q) \bmod q = x \bmod q$

- For these reasons:

$h(i) = ((t[i]* b\text{^}M\text{-}1 \bmod q) + (t[i+1]* b\text{^}M\text{-}2 \bmod q) + \ldots + (t[i+M-1] \bmod q)) \bmod q$

$h(i+1) = (h(i)* b \bmod q$ (Shift left one digit)

$-t[i]* b\text{^}M \bmod q$ (Subtract leftmost digit)

$+t[i+M] \bmod q$) (Add new rightmost digit)

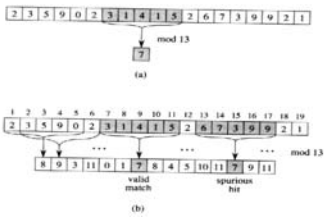
$\bmod q$

Rabin Karp Algorithm

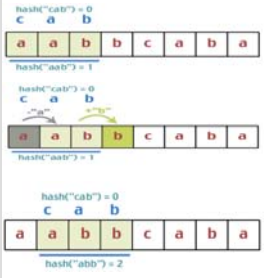
```
1. n = Length[T],m = Length[P]
2. h = d^m-1 mod q
3. p = 0,to = 0
4. for i= 0 to m
5.     do p = (d*p + P[i]) mod q
6.        to = (d*to + T[i]) mod q
7. For s =0 to n-m
8.     do if p=ts
9.           then if P[1..m] ==T[s+1..s+m]
10.                 then “Pattern founds at shift “ s
11.          if s<n-m
12.          then ts+1=(d(ts-T[s+1])h+T[s+m+1]
```

Example of R.K. Algorithm

Calculation for 14152: ts+1=(d(ts-T[s+1]h)+T[s+m+1])mod q
10(31415–3*((10^4)mod13))+2 mod 13
=10(31415- 3*3)+2 mod 13
=8

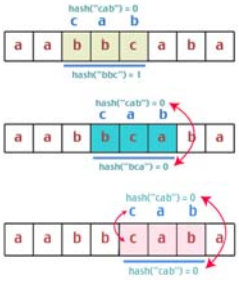


Example



- Let's say our pattern is "cab". And let's say our text string is "aabbcaba".
- For the sake of clarity, we'll use 0 through 26 here to represent letters as opposed to their actual ASCII values.
- For simplicity we will take mod 3.

Example



- Here we have collision but it is ignored by code line 8-10

Running Time Of R.K. Algorithm

- Running time for Rabin Karp algorithm is O((n-m+1)m) in the worst case, since the Rabin Karp algorithm explicitly verifies every valid shift.

Applications

- Text processing
- Bioinformatics
- Compression

References

- Introduction to Algorithm
  -Thomas H. Corman,Ronald L. Rivest, Charles F. Leiserson.
- http://www.sparknotes.com/cs/searching/hashtables/section4.rhtml
- http://www-igm.univ-mlv.fr/~mac/REC/DOC/B5-survey.html
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Questions Please